

Music Processing

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Dynamic Time Warping (DTW)

Solution to Problem 1

$$\text{DTW}(X, Y) = 3, \quad \text{DTW}(X, Z) = 3, \quad \text{DTW}(Y, Z) = 3.$$

Solution to Problem 2

$$C = \begin{pmatrix} 5 & 4 & 4 & 1 \\ 3 & 2 & 2 & 3 \\ 3 & 2 & 2 & 3 \\ 6 & 5 & 5 & 0 \\ 0 & 1 & 1 & 6 \end{pmatrix} \quad D = \begin{pmatrix} 17 & 13 & 13 & 9 \\ 12 & 9 & 9 & 8 \\ 9 & 7 & 7 & 5 \\ 6 & 5 & 6 & 2 \\ 0 & 1 & 2 & 8 \end{pmatrix}$$

An optimal warping path is given by

$$((1, 1), (1, 2), (1, 3), (2, 4), (3, 4), (4, 4), (5, 4)).$$

There is no other optimal warping path.

Solution to Problem 3

Let $X = (x_1, x_2, \dots, x_N)$ and $Y = (y_1, y_2, \dots, y_M)$ be two arbitrary sequences over \mathcal{F} . Furthermore, let $p = (p_1, \dots, p_L)$ with $p_\ell = (n_\ell, m_\ell) \in [1 : N] \times [1 : M]$, $\ell \in [1 : L]$, be a warping path between X and Y with total cost

$$c_p(X, Y) := \sum_{\ell=1}^L c(x_{n_\ell}, y_{m_\ell}).$$

We define a path $q = (q_1, \dots, q_L)$ by $q_\ell := (m_\ell, n_\ell) \in [1 : M] \times [1 : N]$. Obviously, q defines a warping path between Y and X . Furthermore, because of the symmetry of c , one has $c_q(Y, X) = c_p(X, Y)$. It follows that q is an optimal warping path between Y and X if and only if p is an optimal warping path between X and Y . Hence, $\text{DTW}(X, Y) = \text{DTW}(Y, X)$.

Let $\mathcal{F} = \{\alpha, \beta, \gamma\}$ and $c : \mathcal{F} \times \mathcal{F} \rightarrow \mathbb{R}_{\geq 0}$ defined by $c(x, y) := 1 - \delta_{xy}$. Furthermore, let $X = \alpha\beta\gamma$, $Y = \alpha\beta\beta\gamma$ and $Z = \alpha\gamma\gamma$. Then $\text{DTW}(X, Y) = \text{DTW}(Y, X) = 0$, $\text{DTW}(X, Z) = 1$ and

$$\text{DTW}(Y, Z) = 2 > 1 = \text{DTW}(Y, X) + \text{DTW}(X, Z).$$

Hence, the triangular inequality is violated.

Solution to Problem 4

Let L_k be the length of the warping path at level k between the sequences X_k and Y_k each having length 2^{n-k+1} , $1 \leq k \leq n$. Then,

$$L_k \leq 2 \cdot 2^{n-k+1} = 2^{n-k+2},$$

i. e., $L_1 \leq 2N$, $L_2 \leq 2(N/2) = N$, and so on. Then the following holds:

$$\begin{aligned} A^{\text{MsDTW}}(N) &= A^{\text{MsDTW}}\left(\frac{N}{2}\right) + f_1^2 \cdot L_2 \\ &\leq A^{\text{MsDTW}}\left(\frac{N}{2}\right) + 4 \cdot N \\ &= A^{\text{MsDTW}}\left(\frac{N}{2^2}\right) + f_2^2 \cdot L_3 + 4 \cdot N \\ &\leq A^{\text{MsDTW}}\left(\frac{N}{2^2}\right) + 4 \cdot \left(\frac{N}{2}\right) + 4 \cdot N \\ &\leq \dots \\ &\leq A^{\text{MsDTW}}\left(\frac{N}{2^{n-1}}\right) + 4 \cdot \left(\frac{N}{2^{n-2}}\right) \dots + 4 \cdot \left(\frac{N}{2^2}\right) + 4 \cdot \left(\frac{N}{2}\right) + 4 \cdot N \\ &\leq A^{\text{DTW}}(2) + 4 \cdot (4 + \dots + 2^{n-2} + 2^{n-1} + 2^n) \\ &\leq 4 \cdot \sum_{k=0}^n 2^k \\ &\leq 4 \cdot 2^{n+1} \\ &= 8N \end{aligned}$$

The last inequality can be shown easily via induction.